Hybrid Encryption and Further Topics Cryptography (lecture portion)

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Recap

- Symmetric cryptography
 - Key exchange
 - DH(E)
 - AES
- Asymmetric cryptography
 - Public key
 - Private/secret key
 - RSA
- Cryptographic hashes
 - Collision resistance
 - (Second) pre-image resistance
 - SHA-2

Overview

- Hybrid encryption
 - Key exchange revisited (with key derivation)
 - Hybrid encryption scheme
 - Challenge-response scheme
- Examples of combined use of cryptography
 - Secure password storage
 - Transport Layer Security
 - Practical protocols (choice)

Key exchange revisited I

- Diffie-Hellman key exchange
 - Alice and Bob exponentiate random numbers in a modulus (multiplicative group)
 - Rules for exponents give Alice and Bob the same group element
 - Final group element can be used to derive a shared key
 - Eve cannot feasibly compute the shared key (discrete logarithm problem and Diffie-Hellman assumptions)
- Challenges
 - Choosing safe groups
 - \bullet Converting the final group element into a key \to key derivation

Key exchange revisited II

- Key derivation
 - Derive a key k from the final group element f
 - Example: Get 128-bit AES key from 1,024-bit group element
 - Multiple possibilities
- Hashing for key derivation
 - Compute a cryptographic hash: k := H(f)
 - \rightarrow Removes any remaining algebraic structure
 - → Yields a fixed-size result (key)
- Standardized advanced versions, e.g., PBKDF2
 - Repeated use of hashes, e.g., 1,000 iterations
 - → Slows down attackers significantly
 - Use of additional salt (details later)

[1] Kaliski, B.: PKCS #5: Password-Based Cryptography Specification Version 2.0. https://www.ietf.org/rfc/rfc2898.txt (accessed on October 8, 2022), 2000.

Hybrid encryption I

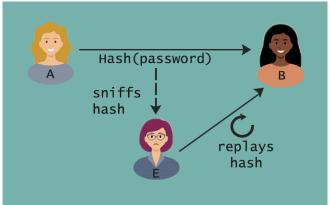
- Alternative to Diffie-Hellman key exchange
- Combine asymmetric and symmetric ciphers
- Assumptions
 - Alice has a public-private key pair $\left(k_{public}^A, k_{private}^A\right)$
 - Bob knows Alice's public key (recall certificates!)
- Key exchange steps:
 - **1** Alice generates a random key k
 - ② Alice encrypts k with her (asymmetric) secret key: $k_e = E(k_{private}^A, k)$
 - **3** Alice sends the encrypted key k_e to Bob
 - Bob decrypts the received key with Alice's public key: $k = D(k_{public}^A, k_e)$
 - Alice and Bob can exchange messages with symmetric key k

Hybrid encryption II

- Why not use asymmetric cryptography for all messages?
 - Public key cryptography is much less efficient than secret key cryptography (longer messages require much more expensive computations)
 - RSA keys are much longer than AES keys for the same level of security (recall elliptic curves for a viable alternative!)
- Summary of hybrid encryption
 - Use public key cryptography to exchange the key
 - Use secret key cryptography to exchange messages
 - The symmetric key is practically random
 - Alice's public key must be known in advance for this to work
- Supplement to this scheme: Challenge-response scheme

Challenge and response I

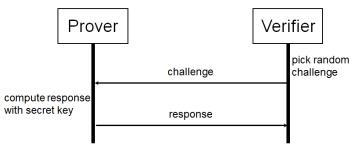
- How to verify that Bob is talking to Alice and not to Eve?
- Risk of replay attacks by Eve



Source: Gibson, D.: Replay Attacks. https://cybersecurityglossary.com/replay-attacks/ (accessed on October 8, 2022), 2020.

Challenge and response II

- How to solve authentication?
 - Certificates (lack liveness, allow for replay attacks)
 - Challenge and response (impedes replay attacks by testing liveness)



Source: Rathgeb, C.: IT-Sicherheit, Kapitel 6 - Authentifikation.

 $\label{lem:https://www.dasec.h-da.de/wp-content/uploads/2014/05/06_authentifikation.pdf (accessed on October 8, 2022), \\ 2014.$

[2] Barrett, D. J. and Silverman, R. E.: SSH: The Secure Shell: The Definitive Guide, 1st ed., O'Reilly, 2001.



Challenge and response III

- Example: Challenge-response scheme for public-key authentication in the Secure Shell (SSH) protocol (simplified):
 - **1** Alice sends her public key k_{public} to Bob
 - ② Bob generates a random challenge r and encrypts it with Alice's supposed public key: $c := E(k_{public}, r)$
 - 3 Bob sends the challenge c to Alice
 - 4 Alice decrypts the supposed challenge c' with her secret key k_{secret} : $r' = D(k_{secret}, c')$
 - **5** Alice hashes the decrypted challenge r': h' := H(r')
 - **1** Alice sends the hashed response h' to Bob
 - **1** Bob hashes its random challenge r: h := H(r)
 - **3** Bob checks whether h = h' (if not, he may be talking to Eve)
- Randomness impedes replay attacks
- Hashing is required so that chosen-plaintext attacks are avoided

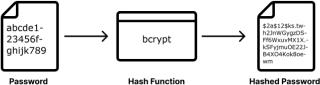
[2] Barrett, D. J. and Silverman, R. E.: SSH: The Secure Shell: The Definitive Guide, 1st ed., O'Reilly, 2001.



Overview of secure password storage

- How to store (user) passwords?
 - Not as plaintext (obvious)
 - Encrypted? → Decryption key yields plaintext passwords
 - Hashed? → Identical passwords yield identical hashes
 - In practice: Salted and stretched





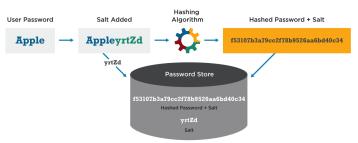
Source: Authgear: Password Hashing and Salting Explained. https://www.authgear.com/post/password-hashing-salting (accessed on October 8, 2022), 2022.

[3] Arias, D.: Adding Salt to Hashing: A Better Way to Store Passwords. https://auth0.com/blog/adding-salt-to-hashing-a-better-way-to-store-passwords/ (accessed on October 8, 2022), 2021.

Salting

- Salt: Random number (later stored together with password) per user
- Salting: Concatenating the password p with a salt s before hashing
- Salted hashing: $\hat{H}(p,s) := H(p||s)$

Password Hash Salting



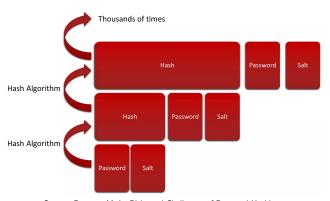
Source: Taylor, C.: Password Salting. https://cyberhoot.com/cybrary/password-salting/ (accessed on October 8, 2022), 2020.

[3] Arias, D.: Adding Salt to Hashing: A Better Way to Store Passwords.

https://auth0.com/blog/adding-salt-to-hashing-a-better-way-to-store-passwords/ (accessed on October 8, 2022),
2021.

Stretching

- Stretching: Prolonging computation by repeated hashing
- Iterations: $P_1 = \hat{H}(p, s), P_{n>1} = \hat{H}(P_{n-1}||p, s)$
- Stored password hash: P_N for large N, e.g., 10,000



Source: Romero, M. I.: Risks and Challenges of Password Hashing. https://www.sitepoint.com/risks-challenges-password-hashing/ (accessed on October 8, 2022), 2014.

Attacks on securely stored passwords I

- Rainbow tables (pre-computed hashes when no salts are used)
- Dictionary attack (try a list of common passwords)
- Brute force (try all passwords from a defined set)



Adopted from Arias, D.: Adding Salt to Hashing: A Better Way to Store Passwords. https://auth0.com/blog/adding-salt-to-hashing-a-better-way-to-store-passwords/ (accessed on October 8, 2022), 2021.

[3] Arias, D.: Adding Salt to Hashing: A Better Way to Store Passwords. https://auth0.com/blog/adding-salt-to-hashing-a-better-way-to-store-passwords/ (accessed on October 8, 2022), 2021.

Attacks on securely stored passwords II

- Brute force attacks depend on the assumed password alphabet
- Example: 8 character-passwords with English alphabet (only upper and lower case): Key space is $52^8 > 50^8 = \left(\frac{100}{2}\right)^8 = \frac{(10^2)^8}{2^8} = \frac{10^{16}}{2^8} = 2^2 \cdot \frac{10^{16}}{2^{10}} \approx 4 \cdot \frac{10^{16}}{10^3} = 4 \cdot 10^{13}$
- Dictionary attacks are only successful when the password is contained in the dictionary (otherwise brute force attack required)
- Stretching can slow down dictionary and brute force attacks
- Salting makes rainbow tables futile
- → Use salting and stretching (better: standardized algorithms)
- → User recommendation: Use a password manager

[3] Arias, D.: Adding Salt to Hashing: A Better Way to Store Passwords. https://auth0.com/blog/adding-salt-to-hashing-a-better-way-to-store-passwords/ (accessed on October 8, 2022), 2021. 4 D > 4 A > 4 B > 4 B > -

Overview of Transport Layer Security (TLS)

- Based on the now-obsolete Secure Sockets Layer (SSL)
- Protocol to provide a secure channel between Alice and Bob
- Guarantees
 - Authentication (Bob is always authenticated, Alice optionally)
 - Confidentiality (only Alice and Bob can read data, but not an attacker)
 - Integrity (data cannot be changed by an attacker without detection)
- Uses a variety of cryptographic primitives, e.g.,
 - Asymmetric ciphers
 - Symmetric ciphers
 - Cryptographic hashes
 - Digital signatures
 - Digital certificates



Recap: Digital signatures

- Generated and added to a message by the signer/sender
- Can be checked by the receiver (and everybody else)
- Preserve message integrity instead of message privacy
- Guarantees
 - The signed message has not been modified
 - The signed message originates from the signer
 - ightarrow Non-repudiation: A signer cannot claim **not** to have signed the message
- Rely on public key cryptography and hashes
 - Signature: Hash of message, encrypted with secret key
 - Signature verification: Compare hash and decrypted signature
 - If the signature check fails, either the message has been tampered with or someone else (a different key) has signed the message

Recap: Digital certificates

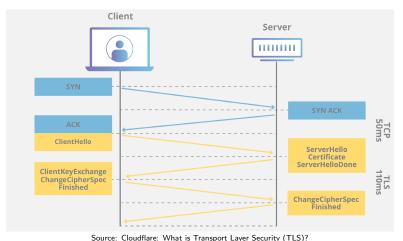
- Digital certificate
 - associates a public key with a person/entity
 - contains information about the signer, the person and their key
 - is signed by a trusted party
 - allows verifying signature and thus identity-to-key matching
 - Assumption: Trusted party whose public key is known by all parties
- Certificate authorities (CAs)
 - are trusted to issue certificates
 - check the identity of the owner or entity
 - are assumed to have a known public key
- Certificate authority hierarchies (CA hierarchies)
 - Issuing of certificates is delegated to sub-CA(s)
 - Multi-level certificates from CA to sub-CA(s) to entity
 - → Hierarchical trust model
 - Root/top-level CA(s) self-sign their own certificate(s)



Key and cipher agreement I

- TLS relies on transport protocols such as TCP
- Initial handshake
 - Authenticates communicating parties through certificates
 - Negotiates used ciphers (not every party supports every cipher)
 - Performs key exchange with forward secrecy
 - Tamper-resistant even against active attackers (who modify messages)
- Used cryptographic primitives
 - Digital certificates for authentication (simplified)
 - Digital signatures for supported ciphers and certificates
 - (EC)DHE and similar protocols for key exchange
 - Hashes of the transcript (sequence of handshake messages)

Key and cipher agreement II



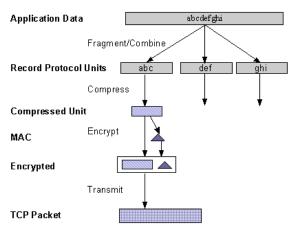
https://www.cloudflare.com/learning/ssl/transport-layer-security-tls/ (accessed on October 26, 2022), 2022.

Secure message exchange I

- TLS can exchange messages after finished handshake
- Record protocol
 - (Alice) Message is split into blocks
 - (Alice) Blocks (records) are protected
 - (Alice) Protected records are transmitted
 - (Bob) Protected records are received
 - (Bob) Protected records are checked and unprotected
 - 6 (Bob) Blocks (records) are reassembled into message
- Used cryptographic primitives
 - Symmetric block ciphers in selected modes (without details)
 - Message authentication codes (MACs, similar to hashes, no details)
- Optional compression step (without details)



Secure message exchange II



Source: The Apache Software Foundation: SSL/TLS Strong Encryption: An Introduction. https://httpd.apache.org/docs/2.4/ssl/ssl_intro.html (accessed on October 26, 2022), 2022.

[4] Rescorla, E.: The Transport Layer Security (TLS) Protocol Version 1.3 (RFC 8446). https://www.rfc-editor.org/rfc/rfc8446.html (accessed on October 26, 2022), 2018.

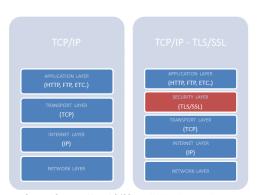
[4]

Practical concerns I

- Newer versions of TLS implement new ciphers and features
- SSL and older TLS versions are
 - vulnerable to certain attacks (without details)
 - support broken ciphers/hash algorithms etc.
 - support weak algorithms and key lengths
- → Deprecation or removal of flawed/weak algorithms and protocol parts
- → Versioning (current: TLS 1.3)
 - Byte-by-byte illustration of a TLS 1.3 connection (including handshake): https://tls13.xargs.org/

Practical concerns II

TLS is designed to provide extra layer for application-level protocols

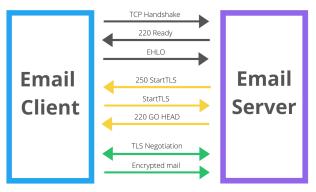


Source: Slaviero, M.: TLS/SSL and .NET Framework 4.0.

https://www.red-gate.com/simple-talk/development/dotnet-development/tlsssl-and-net-framework-4-0/ (accessed on October 26, 2022), 2011.

Practical concerns III

Retrofitting unprotected protocols through STARTTLS command



Source: Griffin, J.: What is StartTLS? https://sendgrid.com/blog/what-is-starttls/ (accessed on October 26, 2022), 2020.

[5] Griffin, J.: What is StartTLS? https://sendgrid.com/blog/what-is-starttls/ (accessed on October 26, 2022), 2020.

Practical protocols (choice)

- Practical cryptographic protocols to illustrate combined use of cryptography (choose one):
 - Masking (smart meter data aggregation) [6]
 - Proof of work blockchains (electronic cash) [7]
 - Wi-Fi Protected Access 3 (WPA3, wireless communication) [8, 9]
 - Content protection (video streaming) [10, 11]
 - Further topics (suggestions)
- [6] Kursawe, K., Danezis, G. and Kohlweiss, M.: Privacy-friendly Aggregation for the Smart Grid. In PETS 2011: Privacy Enhanced Technology Symposium, pp. 175–191, 2011.
- [7] Nakamoto, S.: Bitcoin: A Peer-to-Peer Electronic Cash System. https://bitcoin.org/bitcoin.pdf (accessed on October 26, 2022), 2008.
- [8] WiFi Alliance: WPA3™ Specification Version 3.0.
- https://www.wi-fi.org/downloads-public/WPA3_Specification_v3.0.pdf/35332 (accessed on October 26, 2022), 2020.

 [9] WiFi Alliance: Wi-Fi Protected Access(R) Security Considerations.
- https://www.wi-fi.org/downloads-public/Security_Considerations_20210511.pdf/36067 (accessed on October 26, 2022), 2021.
- [10] Éuropean Telecommunications Standards Institute and European Broadcasting Union: Digital Video Broadcasting (DVB); Content Protection and Copy Management (DVB-CPCM); Part 2: CPCM Reference Model (ETSI TS 102 825-2 V1.2.1), 2011. [11] European Telecommunications Standards Institute and European Broadcasting Union: Digital Video Broadcasting (DVB); Content Protection and Copy Management (DVB-CPCM); Part 5: CPCM Security Toolbox (ETSI TS 102 825-5 V1.2.1), 2011.

Thank you for your attention!

Questions?